

Desnos Anthony (ESIEA SI&S)

« virus don't harm, ignorance does » herm1t





- Outline
 - What's a k-ary virus ???
 - Implementation
 - Conclusion



- What's a k-ary virus ?
 - Cohen's general model of computer viruses :
 - every code is made of a single program which contains the whole instructions devoted to its action



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- What's a k-ary virus ?
 - Cohen's general model of computer viruses :
 - every code is made of a single program which contains the whole instructions devoted to its action

Since every virus is supposed to be composed of a single code, antiviral detection itself considers only this model





- Scattered the viral information over differents files
 - make the viral detection far more complex



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- Scattered the viral information over differents files
 - make the viral detection far more complex

The k constituting part looks like an innocent file and thus does not trigger any alert



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• Definition (Éric Filiol) :

- A k-ary virus is a family of k files (some of them may be not executable) whose union constitues a computer virus and performs an offensive action that is equivalent to that of a true virus. Such a code is said sequential (serial mode) if the k constituent parts are acting strictly one after the another. It is said parallel if the k parts executes simultaneously (parallel mode).



- Two modes :
 - Class I (sequential)
 - codes are executing one after another
 - Class II (parallel)
 - codes are executing at the same time



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With 3 subclasses

- A subclass (dependent sequential codes)
 - « Every part refers or contains a reference to the other ones. It is the weakest class in termof detection since successful detection of one part helps to detect the others. »
- B subclass (independent sequential codes)
 - « No part is referring to another one. Detecting one part does not endanger the other ones. The detected part may be automatically replaced under a different form. »
- C subclass (weakly dependent sequential codes)
 - « Dependency between codes is partial and directed only. »



- Class I (C and B subclasses)
 - The most interesting
- We must make N exploitations to execute the real virus
- Split our virus in differents parts :
 - the first contains the encrypted viral payload
 - the others contain the secret key
- Linux system



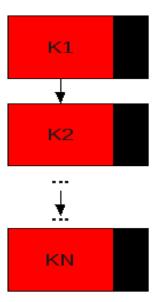
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Main virus

Decryption routine (ASM)

Payload ciphered (ASM + Python)

Others (K virus)





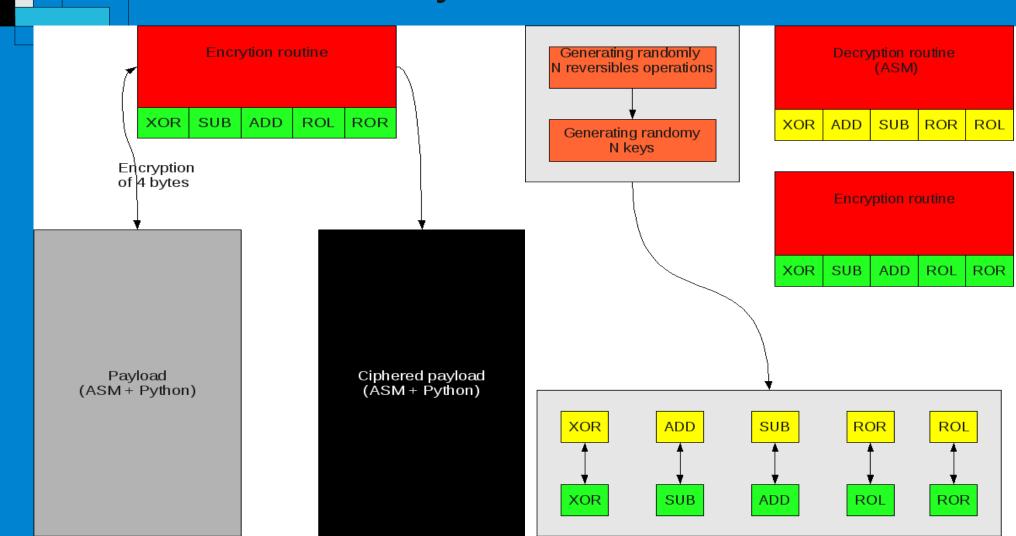
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Six steps:

- Generation of N separate entities, a main entity containing the viral payload (with or without information about secret key), and secondary entities which reconstruct the private key to activate the viral payload,
- The decryption routine,
- Loading of the python script through several techniques,
- Executing of python program, which decrypts with the help of others viruses the final payload,
- Loading of the decrypted payload which is in memory,
- The spread of the virus, in particular the generation of a new routine of encryption and decryption, therefore, with a return to stage 1.

- Polymorphic engine
 - CLET Team (Polymorphic Shellcode Engine Using Spectrum Analysis, Phrack Magazine 61, 2004)
 - Generate a ciphered code which is different at each generation, with different keys
 - generate N reversible operations with N keys
 - examples (simple operations):
 - $XOR \rightarrow XOR$
 - $ADD \rightarrow SUB$
 - $ROL \rightarrow ROR$





- Loading the script
 - Contains a simple script (in python)
 - In this script we have a buffer (or this script can download a buffer ...)
 - which decrypt the final payload
 - when the key is complete
 - How can I execute (stealth) my python script ??

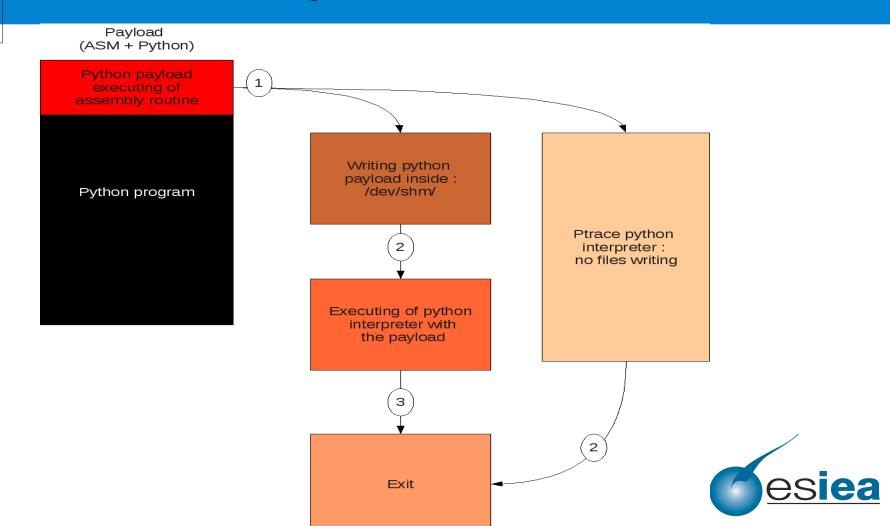


- Loading the script
 - It's written in assembly language
 - We can use /dev/shm
 - tmpfs --> ramfs
 - It's a memory file system!



- Loading the script
 - It's written in assembly language
 - We can use ptrace
 - Hijack open/read/close to load our own code which is in our memory!





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- Remote loading of python code
 - We have download a remote python code which can be:
 - in memory in the same process,
 - in memory in another process,
 - on internet, for example on pastebin.com

How can I execute a remote python code ??



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- Remote loading of python code
 - a simple python class LoadingRemoteModule
 - which gets the buffer, creates classes and calls functions
 - We can use python module:
 - « new » module : creation of run time internal objects
 - with « module » function

➤ mod = new.module (name)



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- Remote loading of python code
 - a simple python class LoadingRemoteModule
 - which gets the buffer, creates classes and calls functions
 - We can use python module :
 - « exec » module : which load a string (or an object of type file, or object code) in a context. This context should be the dictionary of our new module.

exec source in mod.__dict__



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- Remote loading of python code
 - a simple python class LoadingRemoteModule
 - which gets the buffer, creates classes and calls functions
 - We can use python module:
 - Once the module is in the context, it must be load
 - « __import___ » function : returns the module

module = __import__(modulename)



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- Remote loading of python code
 - a simple python class LoadingRemoteModule
 - which gets the buffer, creates classes and calls functions
 - « getattr » function : permits from the module to retrieve a class

class = getattr (module , classname)



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- Remote loading of python code
 - a simple python class LoadingRemoteModule
 - which gets the buffer, creates classes and calls functions
 - « inspect » module :
 - « getargspec » function : to know for a function (thus, the case of constructor) the number of argument, the names and default values.

arg = inspect.getargspec (class.__init__)[0]



- Remote loading of python code
 - a simple python class LoadingRemoteModule
 - which gets the buffer, creates classes and calls functions
 - Then the object is simply constructed with the class returned by getattr and the arguments are in parameters.

```
newinit = []
arg.pop(0)
for i in arg :
    newinit.append(i)

    newargs = izip ( newinit, args )
    d = {}
    for i in newargs :
        d[str (i[0])] = i[1]
    obi = class(d)
```



- Cryptographic library
 - several problems for a virus
 - to use a weak encryption,
 - embedded a tested library or its own optimized library (risk of a poor implementation),
 - to use a library on the system.



- Cryptographic library
 - several problems for a virus
 - to use a weak encryption,
 - embedded a tested library or its own optimized library (risk of a poor implementation),
 - to use a library on the system.
 - We have made the choice to use a library on the system,
 - and therefore take full advantage of a variable present in a vast majority of Linux machines.



- Cryptographic library
 - several problems for a virus
 - to use a weak encryption,
 - embedded a tested library or its own optimized library (risk of a poor implementation),
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- Cryptographic library
 - Openssl in python ?
 - Not in the default python installation



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CTYPES \o/





- Cryptographic library
 - Ctypes
 - Load a dynamic library
 - call its functions



- Cryptographic library
 - Ctypes
 - Load libssl

```
OPENSSL_FILENAME = find_library ("ssl")

→ openssl = cdll.LoadLibrary(OPENSSL FILENAME)
```



- Cryptographic library
 - Ctypes
 - RSA: generate new pairs of key

```
openssl.RAND_load_file ("/dev/random", 2048)
    rsa = c_void_p (openssl.RSA_generate_k e y (bits, 0x10001, None, None))
    rsa_size = openssl.RSA_size (rsa.value)
```



- Cryptographic library
 - Ctypes
 - RSA: Encrypt/Decrypt

```
o = create_string_buffer(rsa_size)
input = create_string_buffer(buffer[i:i+self.rsa_size - 11])
openssl.RSA_public_encrypt(len(input.raw) - 1, addressof(input), addressof(o), rsa.value, 1)

o = create_string_buffer(rsa_size)
input = create_string_buffer(buffer[i:i+self.rsa_size - 11])
openssl.RSA_public_decrypt(len(input.raw) - 1, addressof(input), addressof(o), rsa.value, 1)
```

- Cryptographic library
 - Ctypes
 - RSA: private key without encryption in the PEM format

```
rsa_private_key = ""
bio = c_void_p(self.openssl.BIO_new(openssl.BIO_s_mem()))
if openssl.PEM_write_bio_RSAPrivateKey(bio.value, rsa.value, None, None) == 1:
    temp = c_char_p()
    bufpriv_len = openssl.BIO_ctrl(bio.value, 3, 0, addressof(temp))
    tmp = temp.value
    rsa_private_key = tmp[0:bufpriv_len]
```

- Cryptographic library
 - Ctypes
 - AES

```
class AES_KEY(Structure):
    _fields_ = (
        ("rd_key", c_uint * 60),
        ("rounds", c_int),
      )

enc_key = AES_KEY()
dec_key = AES_KEY()

openssl.AES_set_encrypt_key(key, 16 * 8, addressof(enc_key))
openssl.AES_set_decrypt_key(key, 16 * 8, addressof(dec_key))
```

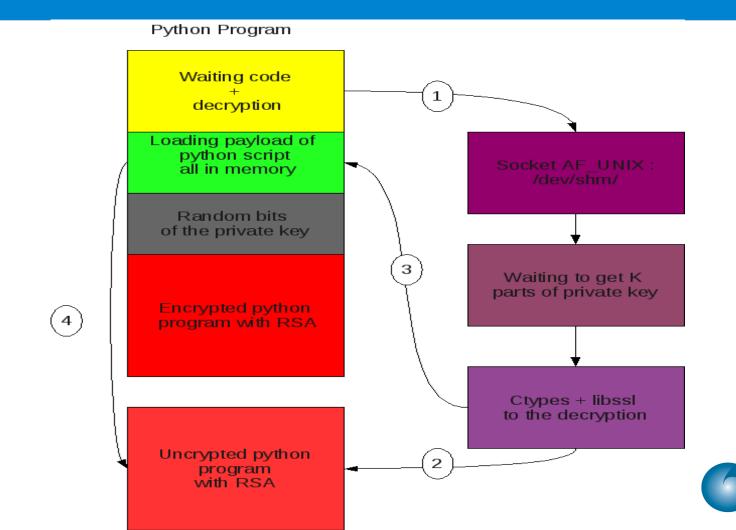


- Cryptographic library
 - Ctypes
 - AES

```
o = create_string_buffer(16)
openssl.AES_encrypt(buffer[i:i+16], addressof(o), addressof(enc_key))
```

```
o2 = create_string_buffer(16)
openssl.AES_decrypt(addressof(o), addressof(o2), addressof(dec_key))
```





- Our main problem is to protect our final payload
 - We have encrypted it, but it remains the problem of the storage of the key
 - If the key is contained in the same source code that the virus,
 then it is very easy for an analyst to find it



- Our main problem is to protect our final payload
 - We have encrypted it, but it remains the problem of the storage of the key
 - If the key is contained in the same source code that the virus, then it is very easy for an analyst to find it
 - K -ary viruses can provide an elegant solution to this problem.



- K-ary virus in sequential mode, C subclass
 - weakly dependent sequential codes
 - split our key in equal parts
 - in some cases that could allow an analyst to have all parts of the key



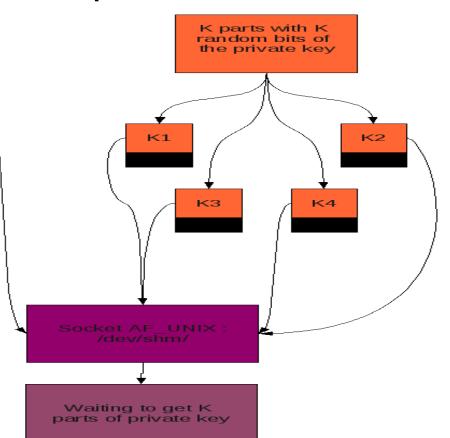
- K-ary virus in sequential mode, C subclass
 - weakly dependent sequential codes
 - split our key in equal parts but also randomly
 - thus it is impossible for an analyst to retrieve the key without having all different codes



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K-ary virus in sequential mode, C subclass

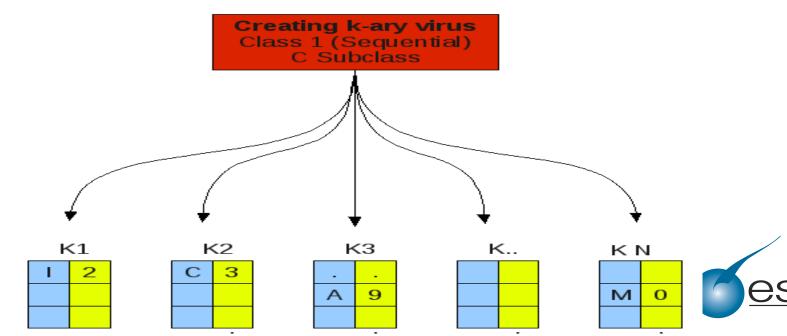






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K-ary virus in sequential mode, C subclass



- K-ary virus in sequential mode, B subclass
 - independent sequential codes
 - the previous subclass has a big flaw, all codes must arrived in the target to start the final payload
 - packets drop
 - missed exploits



- K-ary virus in sequential mode, B subclass
 - independent sequential codes
 - the previous subclass has a big flaw, all codes must arrived in the target to start the final payload
 - packets drop
 - missed exploits
 - it is possible that a code can't arrive and therefore that the spread doesn't continue!



- K-ary virus in sequential mode, B subclass
 - not dependent and can regenerate themselves
 - so if there was a threshold on the different codes generated for the reconstruction of the key without that the totality reaches the destination, or the activation of the final charge after a given time
 - it would continue the spread



- K-ary virus in sequential mode, B subclass
 - not dependent and can regenerate themselves
 - so if there was a threshold on the different codes generated for the reconstruction of the key without that the totality reaches the destination, or the activation of the final charge after a given time
 - it would continue the spread
 - secret share schemes



- K-ary virus in sequential mode, B subclass
 - secret share schemes
 - the main goal is to divide a data D into n pieces
 D1Dn in the following manner between different participants
 - knowledge of any k or more Di pieces makes D easily computable,
 - knowledge of any k 1 or fewer Di pieces leaves D completely undetermined.

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- K-ary virus in sequential mode, B subclass
 - Shamir's secret sharing
 - 2 points are sufficient to define a line,
 - 3 points are sufficient to define a parabola,
 - 4 points are sufficient to define a cubic curve,

• ...



- K-ary virus in sequential mode, B subclass
 - Shamir's secret sharing
 - take k points to define a polynomial of degree k 1
 - To build the polynomial, choose at random (k 1) coefficients , and let be the secret :

$$f(x) = a_0 + a_1 x + a_2 x + a_3 x + \dots + a_{(k-1)} x^{(k-1)}$$



- K-ary virus in sequential mode, B subclass
 - Shamir's secret sharing
 - Every participant (in our case, every virus) is given from a point X of this system, a pair (X, f (X)) (where each X must be different). When k participants are present, the secret can be found, otherwise it is impossible to recover it.



- K-ary virus in sequential mode, B subclass
 - Shamir's secret sharing
 - Our secret is our private key, a simple solution to handle our key is to transform it into PEM format, and convert it into a big integer
 - Another solution isn't to share the private key but the password which encrypt the key, this reduces the computing time and the data exchanges.
 def str2long(s):
 - Python

```
"""Convert a string to a long integer."""

if type(s) not in (types.StringType, types.UnicodeType):
    raise ValueError, 'the input must be a string'

I = 0L

for i in s:
    I <<= 8
    I |= ord(i)
```

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- K-ary virus in sequential mode, B subclass
 - Neville-Aitken's algorithm
 - Once a virus arrived with its pair (X, f (X)), we must be able to find the secret (our a0). To do this we can use Neville-Aitken 's algorithm to find a coefficient, that allows to calculate any degree of the polynomial:

$$p_{(i,i)}(x) = y_i, 0 \le i \le n, p_{(i,j)}(x) = \frac{((x - x_j)p_{(i,j-1)}(x) + (x_i - x)p_{(i+1,j(x))})}{(x_i - x_j)}, 0 \le i < j \le n.$$



- Desnos Anthony (ESIEA SI&S)
- K-ary virus in sequential mode, B subclass
 - Neville-Aitken's algorithm
 - In this case, we want the coefficient of degree 0 (which is the key or the password):

$$p_{(i,i)}(x) = y_i, \ 0 \le i \le n, \ p_{(i,j)}(x) = \frac{((0-x_j) p_{(i,j-1)}(x) + (x_i-0) p_{(i+1,j(x))})}{(x_i-x_j)}, \ 0 \le i < j \le n.$$



- K-ary virus in sequential mode, B subclass
 - Neville-Aitken's algorithm
 - This algorithm has a space and time complexity both in O(n^2), and can be implemented easily in python

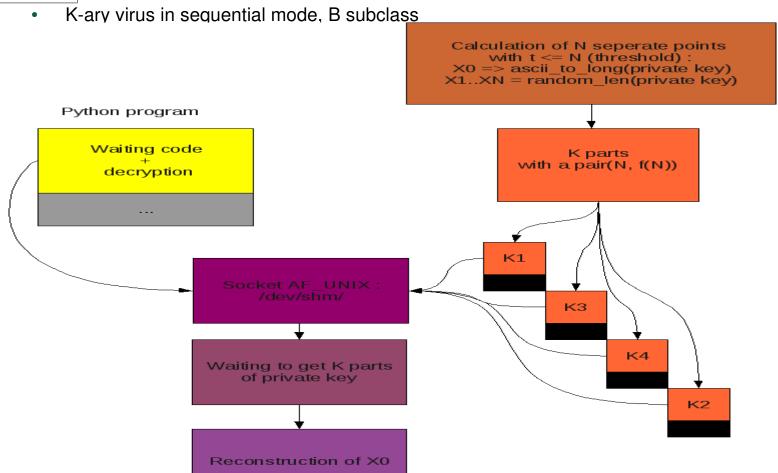
```
def interpolate(x0, y0, x1, y1, x) :
    return (y0*(x-x1) - y1*(x-x0)) / (x0 - x1);

def solveSystem(xs, ys):
    for i in range(1, len(xs)) :
        for k in range(0, len(xs) - i) :
            ys[k] = interpolate(xs[k], ys[k], xs[k+i], ys[k+1], 0)

    return ys[0]
```



```
./shamir.py toto
SECRET toto => TO LONG 1953461359
HASH SECRET
31f7a65e315586ac198bd798b6629ce4903d0899476d5741a9f32e2e521b6a66
f(x) = 1953461359 + 1082694448 x^1 + 100363181 x^2
POINT[1] = 3136518988
POINT[2] = 4520302979
POINT[3] = 6104813332
POINT[4] = 7890050047
POINT[5] = 9876013124
POINT[6] = 12062702563
Running Neville's algorithm: Found x[0]
SECRET = toto
HASH = 31f7a65e315586ac198bd798b6629ce4903d0899476d5741a9f32e2e521b6a66
```

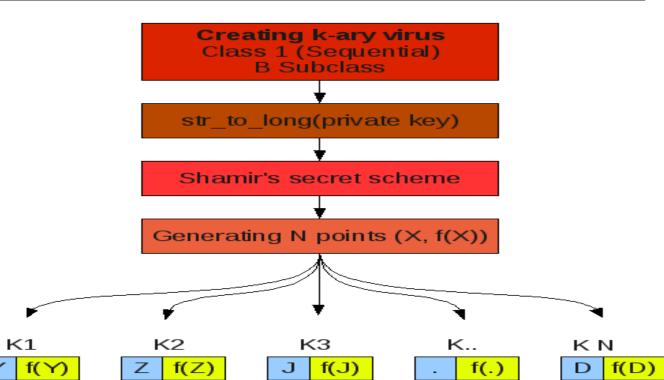




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K-arv virus in sequential mode. B subclass
 Private Key (PEM format)

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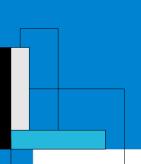


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Conclusion

- K-ary viruses provide an interesting solution to share the key in a virus
- K-ary viruses are a profound change in the way of analysis from the point of view of antivirus





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Many thanks for your attention! Have you any question...?

Happy Hacking!

